

Quantum Stabilizer Codes

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June 10, 2010

A Primer of Group Theory

Groups

A group is a set \mathcal{G} together with a binary operation $*$ on it such that

- **Associativity** : $(f * g) * h = f * (g * h) \forall f, g, h \in \mathcal{G}$;
- **Identity element** : there is an element $e \in \mathcal{G}$ such that

$$g * e = e * g = g;$$

- **Inverse** : for each $g \in \mathcal{G}$, there exists an $h \in \mathcal{G}$ such that

$$g * h = h * g = e.$$

If in addition,

- **Commutativity** : $f * g = g * f \forall f, g \in \mathcal{G}$,

then \mathcal{G} is called an **abelian** group.

Elementary Properties

- The identity element e of a group \mathcal{G} is **unique**.
- The inverse g^{-1} of an element $g \in \mathcal{G}$ is **unique**.

Examples of Groups

- $(C, +)$: the set C of all complex numbers together with addition operation $+$, where 0 is the identity element, is an **abelian** group.
- (C^*, \cdot) : the set C^* of all non-zero complex numbers together with multiplication operation \cdot , where 1 is the identity element, is an **abelian** group.
- $(M_{n \times n}, +)$: the set $M_{n \times n}$ of all $n \times n$ complex matrices together with matrix addition $+$, where the zero matrix $0_{n \times n}$ is the identity element, is an **abelian** group.
- $(M_{n \times n}^*, \cdot)$: the set $M_{n \times n}^*$ of all invertible $n \times n$ complex matrices together with matrix multiplication \cdot , where the identity matrix $I_{n \times n}$ is the identity element, is a **non-abelian** group.

The Pauli Group \mathcal{G}_1 on the State Space of a Qubit

$$\mathcal{G}_1 = \{\pm I, \pm iI, \pm X, \pm iX, \pm Y, \pm iY, \pm Z, \pm iZ\}.$$

- $X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}, Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}.$
- **Hermitian** : $X^\dagger = X, Y^\dagger = Y, Z^\dagger = Z.$
- **Unitary** : $X^2 = Y^2 = Z^2 = I.$
- **Anti-commutative** : $XY = -YX, YZ = -ZY, ZX = -XZ.$
- $XY = iZ, YZ = iX, ZX = iY.$
- \mathcal{G}_1 is non-abelian.

Subgroups of a Group $(\mathcal{G}, *)$

A subset \mathcal{H} of \mathcal{G} is called a subgroup of \mathcal{G} , denoted as $\mathcal{H} < \mathcal{G}$, if

- \mathcal{H} is **closed** under the binary operation $*$, i.e.,

$$f * h \in \mathcal{H}, \forall f, h \in \mathcal{H};$$

- \mathcal{H} is a **group** itself under the binary operation $*$.

Examples of Subgroups

- $(\mathbb{R}, +) < (\mathbb{C}, +)$.
- $(\mathbb{R}^*, \cdot) < (\mathbb{C}^*, \cdot)$.
- $(H_{n \times n}, +) < (M_{n \times n}, +)$, where $H_{n \times n}$ is the set of all Hermitian $n \times n$ complex matrices.
- $(U_{n \times n}, \cdot) < (M_{n \times n}^*, \cdot)$, where $U_{n \times n}$ is the set of all unitary $n \times n$ complex matrices.
- $\mathcal{H} = \{I, X\}$ is a subgroup of the Pauli group \mathcal{G}_1 on a qubit.

The Subgroup Generated by a Set

- $(\mathcal{G}, *)$: a group.
- \mathcal{W} : a subset of \mathcal{G} .
- $\{\mathcal{H}_i, i \in I\}$: the collection of all subgroups of \mathcal{G} which contain \mathcal{W} .
- $\langle \mathcal{W} \rangle \triangleq \bigcap_{i \in I} \mathcal{H}_i$: the subgroup of \mathcal{G} generated by the set \mathcal{W} .
- $\langle \mathcal{W} \rangle = \{g_1^{n_1} * g_2^{n_2} * \cdots * g_k^{n_k}, g_i \in \mathcal{W}, n_i \in \mathbb{Z}, 1 \leq i \leq k, k \geq 1\}$.
- Elements of \mathcal{W} are called the **generators** of the subgroup $\langle \mathcal{W} \rangle$.

Examples

- $\langle 1 \rangle = \{\dots, -2, -1, 0, 1, 2, \dots\} < (C, +)$.
- $\langle 1 \rangle = \{1\} < (C^*, \cdot)$.
- $\langle X \rangle = \{I, X\} < \mathcal{G}_1$.
- $\langle X, Y \rangle = \{\pm I, \pm X, \pm Y, \pm iZ\} < \mathcal{G}_1$.

$$X^{n_1} Y^{n_2} X^{n_3} Y^{n_4} \dots, n_i \in \mathbb{Z}$$

$$= (XY)^n X, (YX)^n Y, (XY)^n, (YX)^n, n \geq 0,$$

$$= \pm X, \pm Y, (\pm iZ)^n, n \geq 0, \text{ since } XY = iZ, YX = -XY,$$

$$(XY)^2 = (YX)^2 = -I, XYX = -Y, YXY = -X,$$

$$= \pm I, \pm X, \pm Y, \pm iZ.$$

- $\langle X, Y, Z \rangle = \{\pm I, \pm iI, \pm X, \pm iX, \pm Y, \pm iY, \pm Z, \pm iZ\} = \mathcal{G}_1$.
 - X, Y, Z are generators of the Pauli group \mathcal{G}_1 .

Independent Generators of a Group

- $\mathcal{H} = \langle g_1, g_2, \dots, g_k \rangle$.
- $\langle g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_k \rangle \neq \mathcal{H}$ for all $i, 1 \leq i \leq k$.
- $\{g_1, g_2, \dots, g_k\}$ is a set of independent generators of \mathcal{H} .

**The Action of the Pauli Group \mathcal{G}_n on the
State Space V of an n -Qubit System**

The State Space V of an n -Qubit System

$$V = V_1 \otimes V_2 \otimes \cdots \otimes V_n.$$

- V_i : the state space of the i th qubit, a two-dimensional complex inner product space with a standard basis $\{|0\rangle, |1\rangle\}$.
- V is a 2^n -dimensional complex inner product space with a standard basis $\{|m\rangle, 0 \leq m \leq 2^n - 1\}$.
 - An example : for $n = 4$,

$$|13\rangle = |1101\rangle = |1\rangle \otimes |1\rangle \otimes |0\rangle \otimes |1\rangle$$

is a state vector in the standard basis.

The Pauli Group \mathcal{G}_n on the State Space V of n Qubits

- $X_i \triangleq \underbrace{I \otimes \cdots \otimes I}_{i-1} \otimes X \otimes \underbrace{I \otimes \cdots \otimes I}_{n-i}$.
- $Y_i \triangleq \underbrace{I \otimes \cdots \otimes I}_{i-1} \otimes Y \otimes \underbrace{I \otimes \cdots \otimes I}_{n-i}$.
- $Z_i \triangleq \underbrace{I \otimes \cdots \otimes I}_{i-1} \otimes Z \otimes \underbrace{I \otimes \cdots \otimes I}_{n-i}$.
- $\mathcal{G}_n \triangleq \langle X_1, Y_1, Z_1, \dots, X_n, Y_n, Z_n \rangle$.
- $i^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n})$: elements in \mathcal{G}_n , where $k_0, k_1, \dots, k_n \in \{0, 1, 2, 3\}$ and $\sigma_0 = I, \sigma_1 = X, \sigma_2 = Y, \sigma_3 = Z$.
- $|\mathcal{G}_n| = 4^{n+1}$.

Properties of \mathcal{G}_n

- Two elements g, h in \mathcal{G}_n is either commutative or anti-commutative, i.e., $gh = \pm hg$.
- Each element g in \mathcal{G}_n is unitary, i.e., $gg^\dagger = g^\dagger g = I$.
- Each element g in \mathcal{G}_n has either $g^2 = I$ or $g^2 = -I$ and then either $g^\dagger = g$ or $g^\dagger = -g$.
- The set of eigenvalues of a $g \neq \pm I, \pm iI$ in \mathcal{G}_n is either $\{1, -1\}$ or $\{i, -i\}$.

Proof

Consider two elements g, h in \mathcal{G}_n such that

$$g = i^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n}), \quad h = i^{m_0} (\sigma_{m_1} \otimes \sigma_{m_2} \otimes \cdots \otimes \sigma_{m_n}).$$

Then we have

$$\begin{aligned} gh &= i^{k_0+m_0} (\sigma_{k_1} \sigma_{m_1} \otimes \sigma_{k_2} \sigma_{m_2} \otimes \cdots \otimes \sigma_{k_n} \sigma_{m_n}) \\ &= (-1)^l i^{k_0+m_0} (\sigma_{m_1} \sigma_{k_1} \otimes \sigma_{m_2} \sigma_{k_2} \otimes \cdots \otimes \sigma_{m_n} \sigma_{k_n}) = \pm hg, \end{aligned}$$

where l is the number of indices i such that $(\sigma_{k_i}, \sigma_{m_i})$ is one of the pairs $(X, Y), (Y, X), (Y, Z), (Z, Y), (Z, X), (X, Z)$. Now $g^\dagger = (-i)^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n})$ and then $gg^\dagger = g^\dagger g = I$, i.e., g is unitary. Also $g^2 = (-1)^{k_0} I = I$ if $k_0 = 0, 2$ and $-I$ if $k_0 = 1, 3$ and then $g^\dagger = g$ if $k_0 = 0, 2$ and $-g$ if $k_0 = 1, 3$. Thus the minimal polynomial of g is a divisor of $x^2 - 1$ or $x^2 + 1$ and if $g \neq \pm I, \pm iI$, the set of eigenvalues of g is either $\{1, -1\}$ or $\{i, -i\}$. \square

Lemma

If \mathcal{H} is a non-abelian subgroup of \mathcal{G}_n , then $-I \in \mathcal{H}$.

Proof. Assume that $-I \notin \mathcal{H}$. Then $g^2 = I$ for all $g \in \mathcal{H}$. Suppose there exists h, u in \mathcal{H} ,

$$h = (-1)^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n}), \quad u = (-1)^{m_0} (\sigma_{m_1} \otimes \sigma_{m_2} \otimes \cdots \otimes \sigma_{m_n}),$$

such that $hu = -uh$. Then we have

$$\begin{aligned} hu &= (-1)^{k_0+m_0} (\sigma_{k_1} \sigma_{m_1} \otimes \sigma_{k_2} \sigma_{m_2} \otimes \cdots \otimes \sigma_{k_n} \sigma_{m_n}) \\ &= (-1)^l (-1)^{m_0+k_0} (\sigma_{m_1} \sigma_{k_1} \otimes \sigma_{m_2} \sigma_{k_2} \otimes \cdots \otimes \sigma_{m_n} \sigma_{k_n}) \\ &= (-1)^l uh, \end{aligned}$$

where l is the number of indices i such that $(\sigma_{k_i}, \sigma_{m_i})$ is one of the pairs $(X, Y), (Y, X), (Y, Z), (Z, Y), (Z, X), (X, Z)$. It is clear that l is an odd number. Note that for each i in the l indices in above,

we have $\sigma_{k_i}\sigma_{m_i}$ to be one of $iZ, -iZ, iX, -iX, iY, -iY$, and for other i , we have $\sigma_{k_i}\sigma_{m_i}$ to be one of I, X, Y, Z . Thus we have

$$\begin{aligned} (hu)^2 &= (-1)^{2(k_0+m_0)}((\sigma_{k_1}\sigma_{m_1})^2 \otimes (\sigma_{k_2}\sigma_{m_2})^2 \otimes \cdots \otimes (\sigma_{k_n}\sigma_{m_n})^2) \\ &= (-1)^l(I \otimes I \otimes \cdots \otimes I) = -I \in \mathcal{H}, \end{aligned}$$

a contradiction. □

Stabilizers under the Action of \mathcal{G}_n on V

- \mathbf{v} : a vector in the state space V of n qubits.
- $\mathcal{G}_n(\mathbf{v}) = \{g \in \mathcal{G}_n \mid g(\mathbf{v}) = \mathbf{v}\}$: the **stabilizer** of \mathbf{v} in \mathcal{G}_n . Note that $\mathcal{G}_n(\mathbf{v}) < \mathcal{G}_n$.
- W' : a set of vectors in V
- $\mathcal{G}_n(W') = \{g \in \mathcal{G}_n \mid g(\mathbf{v}) = \mathbf{v} \forall \mathbf{v} \in W'\} = \bigcap_{\mathbf{v} \in W'} \mathcal{G}_n(\mathbf{v})$: the **stabilizer** of W' in \mathcal{G}_n .
- W : the subspace of V spanned by W' .
- $\mathcal{G}_n(W) = \mathcal{G}_n(W')$: for $g \in \mathcal{G}_n(W')$, we have

$$g\left(\sum_i \alpha_i \mathbf{u}_i\right) = \sum_i \alpha_i g(\mathbf{u}_i) = \sum_i \alpha_i \mathbf{u}_i,$$

where $\alpha_i \in \mathbb{C}$ and $\mathbf{u}_i \in W'$ for all i . Then $g \in \mathcal{G}_n(W)$.

Examples of Stabilizers

- $\mathcal{G}_n(\mathbf{0}) = \mathcal{G}_n$.

- For $n = 2$ and $|\psi\rangle = (|00\rangle + |11\rangle)/\sqrt{2}$, we have

$$\mathcal{G}_2(|\psi\rangle) = \langle X_1 X_2, Z_1 Z_2 \rangle < \mathcal{G}_2.$$

- For $n = 2$ and $|\psi\rangle = (|01\rangle + |10\rangle)/\sqrt{2}$, we have

$$\mathcal{G}_2(|\psi\rangle) = \langle X_1 X_2, -Z_1 Z_2 \rangle < \mathcal{G}_2.$$

- For $n = 2$ and $|\psi\rangle = (|00\rangle - |11\rangle)/\sqrt{2}$, we have

$$\mathcal{G}_2(|\psi\rangle) = \langle -X_1 X_2, Z_1 Z_2 \rangle < \mathcal{G}_2.$$

- For $n = 2$ and $|\psi\rangle = (|01\rangle - |10\rangle)/\sqrt{2}$, we have

$$\mathcal{G}_2(|\psi\rangle) = \langle -X_1 X_2, -Z_1 Z_2 \rangle < \mathcal{G}_2.$$

Lemma

- W : a subspace of V .
- \mathcal{S}' : a subset of \mathcal{G}_n .
- $\langle \mathcal{S}' \rangle$: the subgroup of \mathcal{G}_n generated by \mathcal{S}' .

$$\mathcal{S}' \subseteq \mathcal{G}_n(W) \Leftrightarrow \langle \mathcal{S}' \rangle < \mathcal{G}_n(W).$$

Proof. For all $\mathbf{v} \in W$, we have

$$(g_1^{n_1} g_2^{n_2} \cdots g_k^{n_k})(\mathbf{v}) = \mathbf{v}$$

for any $g_i \in \mathcal{S}'$, $n_i \in \mathbb{Z}$, $1 \leq i \leq k$, and $k \geq 1$. □

Theorem

If W is a non-trivial subspace of V , then the stabilizer $\mathcal{G}_n(W)$ of W in \mathcal{G}_n does not contain $-I$. Thus, $\mathcal{G}_n(W)$ does not contain $\pm iI$ neither and is abelian.

Proof. Since W is non-trivial, there is a non-zero vector \mathbf{v} in W . If $-I \in \mathcal{G}_n(W)$, then $(-I)(\mathbf{v}) = \mathbf{v}$. But $(-I)(\mathbf{v}) = -\mathbf{v}$ so that $\mathbf{v} = -\mathbf{v}$ and then $\mathbf{v} = \mathbf{0}$, a contradiction. \square

Fixed Subspace

- g : an element in \mathcal{G}_n .
- $V(g) \triangleq \{\mathbf{v} \in V \mid g\mathbf{v} = \mathbf{v}\}$: the **fixed subspace** of g in V .
 - If 1 is not an eigenvalue of g , i.e., $g = -I$ or $g^2 = -I$, then $V(g)$ is trivial.
 - If 1 is an eigenvalue of g , i.e., $g \neq -I$ and $g^2 = I$, then $V(g)$ is the eigenspace $E_1(g)$ of g corresponding to eigenvalue 1.
- \mathcal{S}' : a subset of \mathcal{G}_n .
- $V(\mathcal{S}') \triangleq \{\mathbf{v} \in V \mid g\mathbf{v} = \mathbf{v} \ \forall g \in \mathcal{S}'\}$: the **fixed subspace** of \mathcal{S}' in V .
 - $V(\mathcal{S}') = \bigcap_{g \in \mathcal{S}'} V(g)$.
 - If 1 is not an eigenvalue of a g in \mathcal{S}' , then $V(g)$ is trivial.

- $\mathcal{S} = \langle \mathcal{S}' \rangle$: the subgroup of \mathcal{G}_n generated by \mathcal{S}' .
- $V(\mathcal{S}) = V(\mathcal{S}')$: for each $\mathbf{v} \in V(\mathcal{S}')$, we have

$$(g_1^{n_1} g_2^{n_2} \cdots g_k^{n_k})(\mathbf{v}) = \mathbf{v}$$

for any $g_i \in \mathcal{S}'$, $n_i \in \mathbb{Z}$, $1 \leq i \leq k$, and $k \geq 1$. Thus $\mathbf{v} \in V(\mathcal{S})$.

- If $-I \in \mathcal{S}$, then $V(\mathcal{S}) = \{\mathbf{0}\}$.
- If $V(\mathcal{S})$ is non-trivial, then \mathcal{S} is an abelian subgroup of \mathcal{G}_n with $-I \notin \mathcal{S}$.

Theorem

If $\mathcal{S} < \mathcal{G}_n$ and $-I \notin \mathcal{S}$, then

- \mathcal{S} is **abelian**;
- $g^2 = I$ for all $g \in \mathcal{S}$;
- $g^\dagger = g$ for all $g \in \mathcal{S}$;
- the eigenvalues of each $g \neq I$ in \mathcal{S} are **1, -1**.

Proof

- If $g^2 = -I$ for some $g \in \mathcal{S}$, then $-I \in \mathcal{S}$, a contradiction. Thus $g^2 = I$ for all $g \in \mathcal{S}$ and since each element in \mathcal{G}_n is unitary, we have $g^\dagger = g$.
- Since $g^2 = I$, the minimal polynomial of g divides $x^2 - 1$. If $g \neq I$, the eigenvalues of g are 1 and -1. □

The Canonical Representation Theorem

If $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle < \mathcal{G}_n$ and $-I \notin \mathcal{S}$, then each g in \mathcal{S} can be represented as

$$g = g_1^{m_1} g_2^{m_2} \cdots g_k^{m_k},$$

where $m_i \in \{0, 1\}$ for all i , and the representation is **unique** for all $g \in \mathcal{S}$ if and only if g_i 's are **independent generators** of \mathcal{S} .

Proof

Since $-I \notin \mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle$, \mathcal{S} is abelian and $g_i^2 = I$ such that every $g \in \mathcal{S}$ can be represented as

$$g = g_1^{m_1} g_2^{m_2} \cdots g_k^{m_k},$$

where $m_i \in \{0, 1\}$ for all i . Assume that the generators g_1, g_2, \dots, g_k of \mathcal{S} are not independent. Then there exists an i , $1 \leq i \leq k$, such that $g_i \in \langle g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_k \rangle$ and then we have

$$g_i = g_1^{m_1} \cdots g_{i-1}^{m_{i-1}} g_{i+1}^{m_{i+1}} \cdots g_k^{m_k},$$

for $m_i \in \{0, 1\}$. But $g_i = g_i$ and thus g_i has two different representations. Next assume that we have two different representations for an element in \mathcal{S}

$$g_1^{l_1} g_2^{l_2} \cdots g_k^{l_k} = g_1^{m_1} g_2^{m_2} \cdots g_k^{m_k}.$$

Let i be the largest index such that $l_i \neq m_i$. Without loss of generality, we assume $l_1 = 1$ and $m_i = 0$. Thus

$$g_i = g_{i-1}^{l_{i-1}} \cdots g_2^{l_2} g_1^{l_1} g_1^{m_1} g_2^{m_2} \cdots g_{i-1}^{m_{i-1}} \in \langle g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_k \rangle,$$

which says that the generators g_1, g_2, \dots, g_k of \mathcal{S} are not independent. □

The Binary Representation Vector $\varphi(g)$ of $g \in \mathcal{G}_n$

- $g = i^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n})$: an element in the Pauli group \mathcal{G}_n , where $k_i \in \{0, 1, 2, 3\}$, $1 \leq i \leq n$.
- $\varphi(g) = (u_1 u_2 \cdots u_n | v_1 v_2 \cdots v_n)$: a binary $2n$ -tuple corresponding to g such that (note that $Y = iXZ$)

σ_{k_i}	$\sigma_0 = I$	$\sigma_1 = X$	$\sigma_2 = Y$	$\sigma_3 = Z$
u_i	0	1	1	0
v_i	0	0	1	1

- An example : $g = (-i)I \otimes I \otimes X \otimes X \otimes Y \otimes Y \otimes Z \otimes Z$,

$$\varphi(g) = (00111100 | 00001111).$$

The Binary Vector Mapping $g \mapsto \varphi(g)$ Is a Group Homomorphism from \mathcal{G}_n into $(F_2^{2n}, +)$

$$\varphi(gh) = \varphi(g) + \varphi(h).$$

	I	X	Y	Z
I	I	X	Y	Z
X	X	I	iZ	$-iY$
Y	Y	$-iZ$	I	iX
Z	Z	iY	$-iX$	I

↔

$+$	(00)	(10)	(11)	(01)
(00)	(00)	(10)	(11)	(01)
(10)	(10)	(00)	(01)	(11)
(11)	(11)	(01)	(00)	(10)
(01)	(01)	(11)	(10)	(00)

The Binary Vector Mapping $g \mapsto \varphi(g)$ Is a Group Homomorphism from \mathcal{G}_n into F_2^{2n}

- The mapping $g \mapsto \varphi(g)$ is a **group epimorphism** from \mathcal{G}_n onto F_2^{2n} with **kernel**

$$\mathcal{K} = \{I, -I, iI, -iI\}.$$

- $\mathcal{G}_n/\mathcal{K} \cong F_2^{2n}$: a **group isomorphism**.
- The quotient group $\mathcal{G}_n/\mathcal{K}$ is **abelian**.
- The coset \bar{g} of g in the quotient group $\mathcal{G}_n/\mathcal{K}$ is

$$\bar{g} = g\mathcal{K} = \{g, -g, ig, -ig\}.$$

The Binary Vector Mapping $g \mapsto \varphi(g)$ Is a Group Homomorphism from \mathcal{G}_n into F_2^{2n}

- If $\mathcal{S} < \mathcal{G}_n$ such that $-I \notin \mathcal{S}$, then \bar{g}, \bar{h} are different cosets in $\mathcal{G}_n/\mathcal{K}$ for any two different elements g, h in \mathcal{S} .
- If $\mathcal{S} < \mathcal{G}_n$ such that $-I \notin \mathcal{S}$ and \mathcal{S} has k independent generators $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle$, then there are exactly 2^k subgroups $\mathcal{S}' < \mathcal{G}_n$ with $-I \notin \mathcal{S}'$ such that $\varphi(\mathcal{S}') = \varphi(\mathcal{S})$. In particular, we have $\mathcal{S}' = \langle (-1)^{m_1} g_1, (-1)^{m_2} g_2, \dots, (-1)^{m_k} g_k \rangle$ with $m_i \in \{0, 1\} \forall i$.

Proof

- Suppose that $\bar{g} = \bar{h}$ for two different elements g, h in \mathcal{S} . Then we must have $h = -g$ since the eigenvalues of g and h are possibly ± 1 and then $hg = -g^2 = -I \in \mathcal{S}$, a contradiction. Let $\bar{\mathcal{S}} = \{\bar{h} | h \in \mathcal{S}\}$. Then we have $|\bar{\mathcal{S}}| = |\mathcal{S}|$, $\bar{\mathcal{S}} < \mathcal{G}_n/\mathcal{K}$ and that $\bar{\mathcal{S}}$ corresponds to $\varphi(\mathcal{S})$ under the binary vector mapping r .
- Let $\mathcal{S}' < \mathcal{G}_n$ with $-I \notin \mathcal{S}'$ such that $\varphi(\mathcal{S}') = \varphi(\mathcal{S})$. Then we have $\bar{\mathcal{S}}' = \bar{\mathcal{S}}$ and there are $m_i \in \{0, 1\}, \forall i$, such that $\langle (-1)^{m_1} g_1, (-1)^{m_2} g_2, \dots, (-1)^{m_k} g_k \rangle < \mathcal{S}'$. If there is an i such that $(-1)^{m_i} g_i \in \langle (-1)^{m_1} g_1, \dots, (-1)^{m_{i-1}} g_{i-1}, (-1)^{m_{i+1}} g_{i+1}, \dots, (-1)^{m_k} g_k \rangle$, then either g_i or $-g_i$ is in $\langle g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_k \rangle$. The former case contradicts to the independence of the k generators g_1, g_2, \dots, g_k . For the latter case, $(-g_i)g_i = -I \in \mathcal{S}$, a contradiction, too. Thus $(-1)^{m_1} g_1, (-1)^{m_2} g_2, \dots, (-1)^{m_k} g_k$ are

independent. Since $|\mathcal{S}'| = |\overline{\mathcal{S}'}| = |\overline{\mathcal{S}}| = |\mathcal{S}|$, \mathcal{S}' and \mathcal{S} has the same number, says k , of independent generators. We conclude that $\mathcal{S}' = \langle (-1)^{m_1} g_1, (-1)^{m_2} g_2, \dots, (-1)^{m_k} g_k \rangle$ and there are exactly 2^k such \mathcal{S}' 's. \square

Test of Commutativity

- g, h : two elements in \mathcal{G}_n .
- $\Lambda = \begin{bmatrix} 0_{n \times n} & I_{n \times n} \\ I_{n \times n} & 0_{n \times n} \end{bmatrix}$: a binary $2n \times 2n$ matrix.
- $gh = hg \Leftrightarrow \varphi(g)\Lambda\varphi(h)^t = 0$.
 - $g = i^{k_0}(\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n}) \Leftrightarrow$
 $\varphi(g) = (u_1 u_2 \cdots u_n | v_1 v_2 \cdots v_n)$
 - $h = i^{m_0}(\sigma_{m_1} \otimes \sigma_{m_2} \otimes \cdots \otimes \sigma_{m_n}) \Leftrightarrow$
 $\varphi(h) = (u'_1 u'_2 \cdots u'_n | v'_1 v'_2 \cdots v'_n)$
 - $gh = (-1)^l hg$, where l is the number of indices i such that
 $\sigma_{k_i} \sigma_{m_i} \in \{XY, YX, YZ, ZY, ZX, XZ\} \Leftrightarrow u_i v'_i + v_i u'_i = 1$

Test of Independent Generators

- $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle < \mathcal{G}_n$.
- $-I \notin \mathcal{S}$.

g_1, g_2, \dots, g_k are independent generators of $\mathcal{S} \iff$

$\varphi(g_1), \varphi(g_2), \dots, \varphi(g_k)$ are linearly independent vectors in F_2^{2n} .

Proof

“ \Leftarrow ” Suppose there exists an i , $1 \leq i \leq k$, such that

$g_i \in \langle g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_k \rangle$. Then we have

$g_i = g_1^{m_1} \cdots g_{i-1}^{m_{i-1}} g_{i+1}^{m_{i+1}} \cdots g_k^{m_k}$, $m_j \in \{0, 1\}$, by the canonical representation theorem. Thus $\varphi(g_i) = \varphi(g_1^{m_1} \cdots g_{i-1}^{m_{i-1}} g_{i+1}^{m_{i+1}} \cdots g_k^{m_k})$

and $\sum_{i=1}^k m_i \varphi(g_i) = 0$ with $m_i = 1$, which implies that

$\varphi(g_1), \varphi(g_2), \dots, \varphi(g_k)$ are linearly dependent binary vectors.

“ \Rightarrow ” Suppose that $\varphi(g_1), \varphi(g_2), \dots, \varphi(g_k)$ are linearly dependent in F_2^{2n} , i.e., $\sum_{i=1}^k m_i \varphi(g_i) = 0$, where $m_i \in \{0, 1\}$ for all i and at least one $m_i \neq 0$, which implies that $\varphi(\prod_{i=1}^k g_i^{m_i}) = \varphi(I)$ and then

$\prod_{i=1}^k g_i^{m_i} = i^{m_0} I$. Since $-I \notin \mathcal{S}$, we must have $m_0 = 0$, i.e.,

$\prod_{i=1}^k g_i^{m_i} = I = \prod_{i=1}^k g_i^0$. By the canonical representations theorem,

g_1, g_2, \dots, g_k are dependent generators of \mathcal{S} . \square

Remarks

- $-I$ is independent in \mathcal{G}_n but $\varphi(-I) = 0$ is dependent in F_2^{2n} .
- X_1, Y_1, Z_1 are independent in \mathcal{G}_n but $\varphi(X_1) + \varphi(Y_1) + \varphi(Z_1) = 0$.
- If $g_1, g_2, \dots, g_n \in \mathcal{G}_n$ have $\varphi(g_1), \varphi(g_2), \dots, \varphi(g_n)$ linearly independent in F_2^{2n} , then they are independent in \mathcal{G}_n .

The Check Matrix of a Set of Generators

- $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle < \mathcal{G}_n$.
- The check matrix of a set of generators g_1, g_2, \dots, g_k of \mathcal{S} is a $k \times 2n$ binary matrix

$$H \triangleq \begin{bmatrix} \varphi(g_1) \\ \varphi(g_2) \\ \vdots \\ \varphi(g_k) \end{bmatrix} .$$

An Example

$$g_1 = IIIXXX, \quad g_2 = IXIXIX, \quad g_3 = ZIZIZI,$$
$$g_4 = YYYIII, \quad g_5 = XYXYXY, \quad g_6 = YZYZYZ.$$

$$\left[\begin{array}{cccccc|cccccc} 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{array} \right].$$

Maximum Rank of Check Matrices of Abelian Pauli Subgroups

- $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle < \mathcal{G}_n$.
- H : the $k \times 2n$ check matrix of \mathcal{S} .

If \mathcal{S} is abelian, then $\text{Rank}(H) \leq n$.

Proof. Let $\text{Row}(A)$ be the row space of a matrix A . Since \mathcal{S} is abelian,

$$H\Lambda H^T = H(H\Lambda)^T = 0$$

$$\Rightarrow \text{Row}(H\Lambda) \subseteq (\text{Row}(H))^\perp$$

$$\Rightarrow \dim(\text{Row}(H\Lambda)) \leq \dim(\text{Row}(H)^\perp)$$

$$\Rightarrow \text{Rank}(H) \leq 2n - \text{Rank}(H), \quad \text{since } \Lambda \text{ is non-singular,}$$

$$\Rightarrow \text{Rank}(H) \leq n,$$

which completes the proof. □

A Technical Lemma

- $\mathcal{S} = \langle g_1, g_2, \dots, g_k \rangle < \mathcal{G}_n : g_1, g_2, \dots, g_k$ are independent generators.
- $-I \notin \mathcal{S}$.

For each generator g_i , $1 \leq i \leq k$, there exists an $h_i \in \mathcal{G}_n$ such that

$$h_i g_i h_i^\dagger = -g_i \quad \text{and} \quad h_i g_j h_i^\dagger = g_j, \quad \forall j \neq i,$$

$$(h_i g_i = -g_i h_i \quad \text{and} \quad h_i g_j = g_j h_i, \quad \forall j \neq i.)$$

Proof

Let A be the $k \times 2n$ check matrix of the independent generators g_1, g_2, \dots, g_k of \mathcal{S} . Since A has k linearly independent rows, the binary matrix $A\Lambda$ has a right inverse, says a $2n \times k$ binary matrix $B = [\mathbf{y}_1, \mathbf{y}_2, \dots, \mathbf{y}_k]$, such that $A\Lambda B = I_{k \times k}$. For each i , $1 \leq i \leq k$, select an h_i in \mathcal{G}_n such that $\varphi(h_i) = \mathbf{y}_i^t$. It is clear that $\varphi(g_i)\Lambda\varphi(h_j)^t = \delta_{ij}$, $1 \leq i, j \leq k$, which implies that $h_j g_i = (-1)^{\delta_{ij}} g_i h_j$. □

Remark

The above technical lemma is also true if g_1, g_2, \dots, g_k in \mathcal{G}_n have linearly independent $\varphi(g_1), \varphi(g_2), \dots, \varphi(g_k)$ in F_2^{2n} even g_i 's may not commute with each other.

Lemma

Let P_1 and P_2 be two projectors from a complex inner product space W onto the subspaces W_1 and W_2 of W , respectively. Then P_1 and P_2 commute if and only if P_1P_2 is a projector. In this case, P_1P_2 is the projector onto $W_1 \cap W_2$.

Proof. If P_1 and P_2 commute, then we have

$$\begin{aligned}(P_1P_2)^2 &= P_1P_2P_1P_2 = P_1^2P_2^2 = P_1P_2, \\ (P_1P_2)^\dagger &= P_2^\dagger P_1^\dagger = P_2P_1 = P_1P_2.\end{aligned}$$

Thus P_1P_2 is a projector from W onto $(P_1P_2)(W)$. Conversely, if P_1P_2 is a projector, then $P_1P_2 = (P_1P_2)^\dagger = P_2^\dagger P_1^\dagger = P_2P_1$. Note that $(W_1 \cap W_2) \subseteq (P_1P_2)(W)$ since $(P_1P_2)(W_1 \cap W_2) = W_1 \cap W_2$. Also $(P_1P_2)(W) = (P_2P_1)(W)$ implies $(P_1P_2)(W) \subseteq (W_1 \cap W_2)$. We conclude that $(P_1P_2)(W) = W_1 \cap W_2$. \square

Remark

Since projectors are normal operators, two projectors P_1 and P_2 on a complex inner product space W commutes if and only if they can be unitarily diagonalized simultaneously, i.e., there exists a common orthonormal eigenbasis for P_1 and P_2 .

Dimension Theorem

Let $\mathcal{S} < \mathcal{G}_n$ with $-I \notin \mathcal{S}$. If \mathcal{S} is generated by $n - k$ independent generators, then $V(\mathcal{S})$ is a 2^k -dimensional complex vector space.

Proof

Let \mathcal{S} be generated by $n - k$ independent generators g_1, g_2, \dots, g_{n-k} such that $V(\mathcal{S}) = \bigcap_{j=1}^{n-k} V(g_j)$. Since $-I \notin \mathcal{S}$, the eigenvalues of each g_i are 1 and -1 with corresponding eigenspaces $E_1(g_j) = V(g_j)$ and $E_{-1}(g_j)$, respectively. For $x_j \in \{0, 1\}$, $(I + (-1)^{x_j} g_j)/2$ is the projector from V onto $E_{(-1)^{x_j}}(g_j)$. Now for each binary $(n - k)$ -tuple $\mathbf{x} = (x_1, x_2, \dots, x_{n-k})$, we let

$$P_{\mathbf{x}} = \frac{\prod_{j=1}^{n-k} (I + (-1)^{x_j} g_j)}{2^{n-k}}.$$

Since \mathcal{S} is abelian, $P_{\mathbf{x}}$ is the projector from V onto $\bigcap_{j=1}^{n-k} E_{(-1)^{x_j}}(g_j)$ by the previous lemma. In particular, $P_{\mathbf{0}}(V) = \bigcap_{j=1}^{n-k} E_1(g_j) = V(\mathcal{S})$. It is apparently that $P_{\mathbf{x}}(V)$'s are orthogonal. Let h_i in \mathcal{G}_n , $1 \leq i \leq n - k$, be in the technical lemma such that $h_i g_j = (-1)^{\delta_{ij}} g_j h_i$. And for each binary $(n - k)$ -tuple \mathbf{x} ,

let $h_{\mathbf{x}} = \prod_{i=1}^{n-k} h_i^{x_i}$. We have

$$h_{\mathbf{x}} g_j = \left(\prod_{i=1}^{n-k} h_i^{x_i} \right) g_j = (-1)^{x_j} g_j \left(\prod_{i=1}^{n-k} h_i^{x_i} \right) = (-1)^{x_j} g_j h_{\mathbf{x}}.$$

Since $h_{\mathbf{x}}$ is unitary, we have $h_{\mathbf{x}} g_j h_{\mathbf{x}}^\dagger = (-1)^{x_j} g_j$, which implies $(I + (-1)^{x_j} g_j)/2 = h_{\mathbf{x}}(I + g_j)/2 h_{\mathbf{x}}^\dagger$ and

$$P_{\mathbf{x}} = h_{\mathbf{x}} P_{\mathbf{0}} h_{\mathbf{x}}^\dagger.$$

Thus $P_{\mathbf{x}}(V)$ and $P_{\mathbf{0}}(V)$ have the same dimension. Finally,

$$\sum_{\mathbf{x}} P_{\mathbf{x}} = \prod_{j=1}^{n-k} \frac{(I + g_j) + (I - g_j)}{2} = I,$$

which implies that $V = \sum_{\mathbf{x}} P_{\mathbf{x}}(V)$ is a sum of 2^{n-k} orthogonal subspaces of the same dimension. Thus the dimension of $V(\mathcal{S}) = P_{\mathbf{0}}(V)$ is 2^k . □

Corollaries

- If W is a subspace of V , then

$$W \subseteq V(\mathcal{G}_n(W)).$$

Proof. Assume $W = \text{Span}(W')$, where W' is a basis of W .

Then $\mathcal{G}_n(W) = \mathcal{G}_n(W')$ and

$V(\mathcal{G}_n(W)) = V(\mathcal{G}_n(W')) = V(\cap_{\mathbf{v}' \in W'} \mathcal{G}_n(\mathbf{v}'))$, i.e.,

$V(\mathcal{G}_n(W)) = \{\mathbf{v} \in V \mid g\mathbf{v} = \mathbf{v} \ \forall g \in \cap_{\mathbf{v}' \in W'} \mathcal{G}_n(\mathbf{v}')\}$, which contains all vectors w' in W' . Thus $W \subseteq V(\mathcal{G}_n(W))$. □

- If $\mathcal{S} < \mathcal{G}_n$ with $-I \notin \mathcal{S}$, then

$$\mathcal{S} = \mathcal{G}_n(V(\mathcal{S})).$$

Proof. Assume \mathcal{S} is generated by \mathcal{S}' , a set of $n - k$ independent generators. Then $\dim V(\mathcal{S}) = \dim V(\mathcal{S}') = 2^k$ by the dimension

theorem. And $\mathcal{G}_n(V(\mathcal{S})) = \mathcal{G}_n(V(\mathcal{S}')) = \mathcal{G}_n(\bigcap_{g' \in \mathcal{S}'} V(g'))$, i.e.,

$$\mathcal{G}_n(V(\mathcal{S})) = \{g \in \mathcal{G}_n \mid gv = v \forall v \in \bigcap_{g' \in \mathcal{S}'} V(g')\},$$

which contains all elements g' in \mathcal{S}' . Thus $\mathcal{S} \subseteq \mathcal{G}_n(V(\mathcal{S}))$.

Now suppose $\mathcal{S} \subsetneq \mathcal{G}_n(V(\mathcal{S}))$ so that $\mathcal{G}_n(V(\mathcal{S}))$ has $n - k + m$ independent generators, $m \geq 1$. Since

$$V(\mathcal{S}) \subseteq V(\mathcal{G}_n(V(\mathcal{S}))),$$

we have

$$\dim(V(\mathcal{S})) = 2^k \leq \dim(V(\mathcal{G}_n(V(\mathcal{S})))) = 2^{k-m},$$

a contradiction. We conclude that $\mathcal{S} = \mathcal{G}_n(V(\mathcal{S}))$. □

Remark

- Let $W \subseteq V$ be a non-trivial subspace of V . Then $W = V(\mathcal{G}_n(W))$ if and only if W is the fixed subspace of a certain $\mathcal{S} < \mathcal{G}_n$ with $-I \notin \mathcal{S}$.

Proof. (\Rightarrow) Notice that $\mathcal{G}_n(W) < \mathcal{G}_n$ and $-I \notin \mathcal{G}_n(W)$. By letting $\mathcal{S} = \mathcal{G}_n(W)$, we have $V(\mathcal{S}) = V(\mathcal{G}_n(W)) = W$.

(\Leftarrow) We have $W = V(\mathcal{S})$. Since $\mathcal{S} = \mathcal{G}_n(V(\mathcal{S}))$, $\mathcal{G}_n(W) = \mathcal{S}$. So $V(\mathcal{G}_n(W)) = V(\mathcal{S}) = W$. \square

- $\mathcal{S} < \mathcal{G}_n$ is a stabilizer group if and only if $-I \notin \mathcal{S}$.

Proof. \mathcal{S} is a stabilizer group if and only if $V(\mathcal{S})$ is nontrivial if and only if $\mathcal{G}_n(V(\mathcal{S})) = \mathcal{S}$ does not contain $-I$. \square

Unitary Operators

- u : a unitary operator on the state space V of n qubits.
- g : an element in \mathcal{G}_n .
- $V(g)$: the fixed subspace of g in V .
- $u(V(g)) = V(ugu^\dagger)$: the fixed subspace of ugu^\dagger in V .
 - $\mathbf{v} \in V(ugu^\dagger) \Leftrightarrow (ugu^\dagger)(\mathbf{v}) = \mathbf{v} \Leftrightarrow g(u^\dagger(\mathbf{v})) = u^\dagger(\mathbf{v}) \Leftrightarrow u^\dagger(\mathbf{v}) \in V(g) \Leftrightarrow \mathbf{v} \in u(V(g))$.
- $\mathcal{S} < \mathcal{G}_n$: a subgroup of \mathcal{G}_n .
- $V(\mathcal{S})$: the fixed subspace of \mathcal{S} in V .
- $u(V(\mathcal{S})) = V(u\mathcal{S}u^\dagger)$: the fixed subspace of $u\mathcal{S}u^\dagger$ in V .

Stabilizer Codes

- $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle < \mathcal{G}_n$: an abelian subgroup of \mathcal{G}_n with $-I \notin \mathcal{S}$ which is generated by $n - k$ independent generators.
- $\mathcal{C}(\mathcal{S})$: an $[n, k]$ stabilizer code which is just the fixed subspace of \mathcal{S} in the state space V of n qubits.

Centralizer and Normalizer

- $\mathcal{S} < \mathcal{G}_n$.
- $C_{\mathcal{G}_n}(\mathcal{S}) \triangleq \{g \in \mathcal{G}_n \mid ghg^\dagger = h \ \forall h \in \mathcal{S}\}$: the centralizer of \mathcal{S} in \mathcal{G}_n .
 - If $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle$, then $g \in C_{\mathcal{G}_n}(\mathcal{S}) \Leftrightarrow gg_i = g_i g$ for all i .
- $N_{\mathcal{G}_n}(\mathcal{S}) \triangleq \{g \in \mathcal{G}_n \mid ghg^\dagger \in \mathcal{S} \ \forall h \in \mathcal{S}\}$: the normalizer of \mathcal{S} in \mathcal{G}_n .
 - If $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle$, then $g \in N_{\mathcal{G}_n}(\mathcal{S}) \Leftrightarrow gg_i g^\dagger \in \mathcal{S}$ for all i .
- $\mathcal{S} \triangleleft N_{\mathcal{G}_n}(\mathcal{S})$ and $N_{\mathcal{G}_n}(\mathcal{S})$ is the largest subgroup of \mathcal{G}_n which contains \mathcal{S} as a normal group.
- $\mathcal{S} \triangleleft \mathcal{G}_n$ if and only if $N_{\mathcal{G}_n}(\mathcal{S}) = \mathcal{G}_n$.

- $C_{\mathcal{G}_n}(\mathcal{S}) < N_{\mathcal{G}_n}(\mathcal{S})$.
 - If $-I \notin \mathcal{S}$, then $C_{\mathcal{G}_n}(\mathcal{S}) = N_{\mathcal{G}_n}(\mathcal{S})$.

Proof.

Let $g \in N_{\mathcal{G}_n}(\mathcal{S})$. Suppose that

$$gg_i = -g_i g$$

for some i , where $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle$. Then

$$gg_i g^\dagger = -g_i g g^\dagger = -g_i \in \mathcal{S}$$

so that $g_i(-g_i) = -I \in \mathcal{S}$, a contradiction. □

Error-Correction Conditions for Stabilizer Codes

If

- \mathcal{S} : an abelian subgroup of \mathcal{G}_n with $-I \notin \mathcal{S}$,
- $\mathcal{C}(\mathcal{S})$: the stabilizer code fixed by \mathcal{S} ,
- $\{E_i\}$: a set of operation elements (error patterns) in \mathcal{G}_n such that $E_i^\dagger E_j \notin N_{\mathcal{G}_n}(\mathcal{S}) \setminus \mathcal{S}$ for all i, j ,

then $\{E_i\}$ are correctable error patterns by the code $\mathcal{C}(\mathcal{S})$.

Proof

- $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle$: g_1, g_2, \dots, g_{n-k} are generators of \mathcal{S} .
- $P = \prod_{i=1}^{n-k} (I + g_i) / 2^{n-k}$: the projector from V onto the code space $\mathcal{C}(\mathcal{S})$.
- Two possibilities : either $E_i^\dagger E_j \in \mathcal{S}$ or $E_i^\dagger E_j \in \mathcal{G}_n \setminus N_{\mathcal{G}_n}(\mathcal{S})$.
- Case I : $E_i^\dagger E_j \in \mathcal{S}$.
 - $\mathcal{C}(\mathcal{S}) \subseteq V(E_i^\dagger E_j)$: $V(E_i^\dagger E_j)$ is the fixed subspace of $E_i^\dagger E_j$ in V .
 - $\mathcal{C}(\mathcal{S})^\perp$: the orthogonal complement of $\mathcal{C}(\mathcal{S})$ in V .
 - $\mathbf{v} = \mathbf{u} + \mathbf{u}^\perp$ in V : orthogonal decomposition of \mathbf{v} with $\mathbf{u} \in \mathcal{C}(\mathcal{S})$ and $\mathbf{u}^\perp \in \mathcal{C}(\mathcal{S})^\perp$.
 - $E_i^\dagger E_j P(\mathbf{v}) = (E_i^\dagger E_j)(\mathbf{u}) = \mathbf{u} = P(\mathbf{v})$ for all \mathbf{v} in V , which implies $E_i^\dagger E_j P = P$.

- $PE_i^\dagger E_j P = P^2 = P$ and then $PE_j^\dagger E_i P = P$.
- Case II : $E_i^\dagger E_j \in \mathcal{G}_n \setminus N\mathcal{G}_n(\mathcal{S})$.
 - $E_i^\dagger E_j$ must anti-commute with some generator of \mathcal{S} , says g_1 without loss of generality.
 - $E_i^\dagger E_j P = E_i^\dagger E_j \prod_{i=1}^{n-k} (I + g_i) / 2^{n-k} = (I - g_1) E_i^\dagger E_j \prod_{i=2}^{n-k} (I + g_i) / 2^{n-k}$.
 - $P(I - g_1) = 0$ since $(I + g_1)(I - g_1) = 0$.
 - $PE_i^\dagger E_j P = 0$ and then $PE_j^\dagger E_i P = 0$.
- $PE_i^\dagger E_j P = \alpha_{ij} P$ with $\alpha = [\alpha_{ij}]$ a Hermitian matrix.
- Conclusion : $\{E_i\}$ is a set of correctable error patterns by the stabilizer code $\mathcal{C}(\mathcal{S})$. □

Syndrome Measurements and Syndromes

- $\mathcal{S} = \langle g_1, g_2, \dots, g_{n-k} \rangle$: an abelian subgroup of \mathcal{G}_n with $-I \notin \mathcal{S}$.
- $\mathcal{C}(\mathcal{S})$: the stabilizer code fixed by \mathcal{S} .
- $\{E_i\}$: a set of correctable error patterns in \mathcal{G}_n by $\mathcal{C}(\mathcal{S})$ such that $E_i^\dagger E_j \notin N_{\mathcal{G}_n}(\mathcal{S}) \setminus \mathcal{S}$ for all i, j .
- $g_j^\dagger = g_j$ and $g_j = (I + g_j)/2 - (I - g_j)/2$: g_j can be regarded as a projective measurement $\{(I + g_j)/2, (I - g_j)/2\}$.
- Two possibilities : $E_i g_j = g_j E_i$, i.e., $E_i g_j E_i^\dagger = g_j$ or $E_i g_j = -g_j E_i$, i.e., $E_i g_j E_i^\dagger = -g_j$.
 - Case I : $E_i g_j = g_j E_i$. For each $|\psi\rangle \in \mathcal{C}(\mathcal{S})$,

$$(I + g_j)/2 E_i(|\psi\rangle) = E_i(I + g_j)/2(|\psi\rangle) = E_i(|\psi\rangle).$$

The measurement of g_j on the (un-normalized) possibly

corrupted state $E_i(|\psi\rangle)$ is +1 with probability one. And the post-measurement state remains unchanged.

– Case I : $E_i g_j = -g_j E_i$. For each $|\psi\rangle \in \mathcal{C}(\mathcal{S})$,

$$(I - g_j)/2 E_i(|\psi\rangle) = E_i(I + g_j)/2(|\psi\rangle) = E_i(|\psi\rangle).$$

The measurement of g_j on the (un-normalized) possibly corrupted state $E_i(|\psi\rangle)$ is -1 with probability one. And the post-measurement state remains unchanged.

- g_1, g_2, \dots, g_{n-k} : $n - k$ syndrome measurements
- $\beta_i = (\beta_{i1}, \beta_{i2}, \dots, \beta_{i(n-k)})$: syndrome vector of the error pattern E_i , where

$$E_i g_j E_i^\dagger = \beta_{ij} g_j.$$

Error Detection

If the syndrome vector β is not equal to $(1, 1, \dots, 1)$, then the channel output is a corrupted state and error is detected.

Error Correction

- If there is only one error pattern E_i associated with the syndrome vector β , then apply E_i^\dagger to any of the syndrome measurement output to recover the channel input state.
- If there are two error patterns E_i and $E_{i'}$ associated with the syndrome vector β , then $E_i g_j E_i^\dagger = E_{i'} g_j E_{i'}^\dagger$ for all j and then $(E_{i'}^\dagger E_i) g_j = g_j (E_{i'}^\dagger E_i)$ for all j . Thus $E_{i'}^\dagger E_i$ is in $C_{\mathcal{G}_n}(\mathcal{S}) = N_{\mathcal{G}_n}(\mathcal{S})$. But $E_{i'}^\dagger E_i \notin N_{\mathcal{G}_n}(\mathcal{S}) \setminus \mathcal{S}$, we have $E_{i'}^\dagger E_i \in \mathcal{S}$. Thus we can apply E_i^\dagger or $E_{i'}^\dagger$ to any of the syndrome measurement output to recover the channel input state.

Minimum Distance of a Stabilizer Code

- $E = i^{k_0} (\sigma_{k_1} \otimes \sigma_{k_2} \otimes \cdots \otimes \sigma_{k_n})$: an error pattern in \mathcal{G}_n .
- $w(E)$: the number of non-identity terms σ_{k_i} in the tensor product representation of E .
- \mathcal{S} : an abelian subgroup of \mathcal{G}_n with $-I \notin \mathcal{S}$.
- $\mathcal{C}(\mathcal{S})$: the stabilizer code fixed by \mathcal{S} .
- $d = \min\{w(E) \mid E \in N_{\mathcal{G}_n}(\mathcal{S}) \setminus \mathcal{S}\}$: the **minimum distance** of the stabilizer code $\mathcal{C}(\mathcal{S})$.
- Let $d \geq 2t + 1$. Then all E in \mathcal{G}_n with $w(E) \leq t$ are correctable error patterns since

$$w(E_i^\dagger E_j) \leq w(E_i) + w(E_j) \leq 2t$$

for any E_i, E_j in \mathcal{G}_n with $w(E_i) \leq t$ and $w(E_j) \leq t$ and then $E_i^\dagger E_j \notin N_{\mathcal{G}_n}(\mathcal{S}) \setminus \mathcal{S}$.